

The realization view

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Abstract

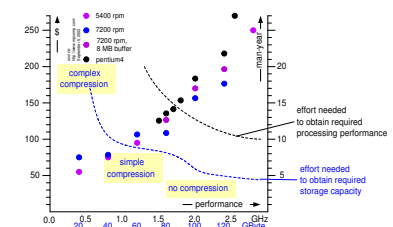
The realization view looks at the actual technologies used and the actual implementation. Methods used here are logarithmic views, micro-benchmarks and budgets.

Analysis methods with respect to safety, reliability and security provide a link back to the functional and conceptual views.

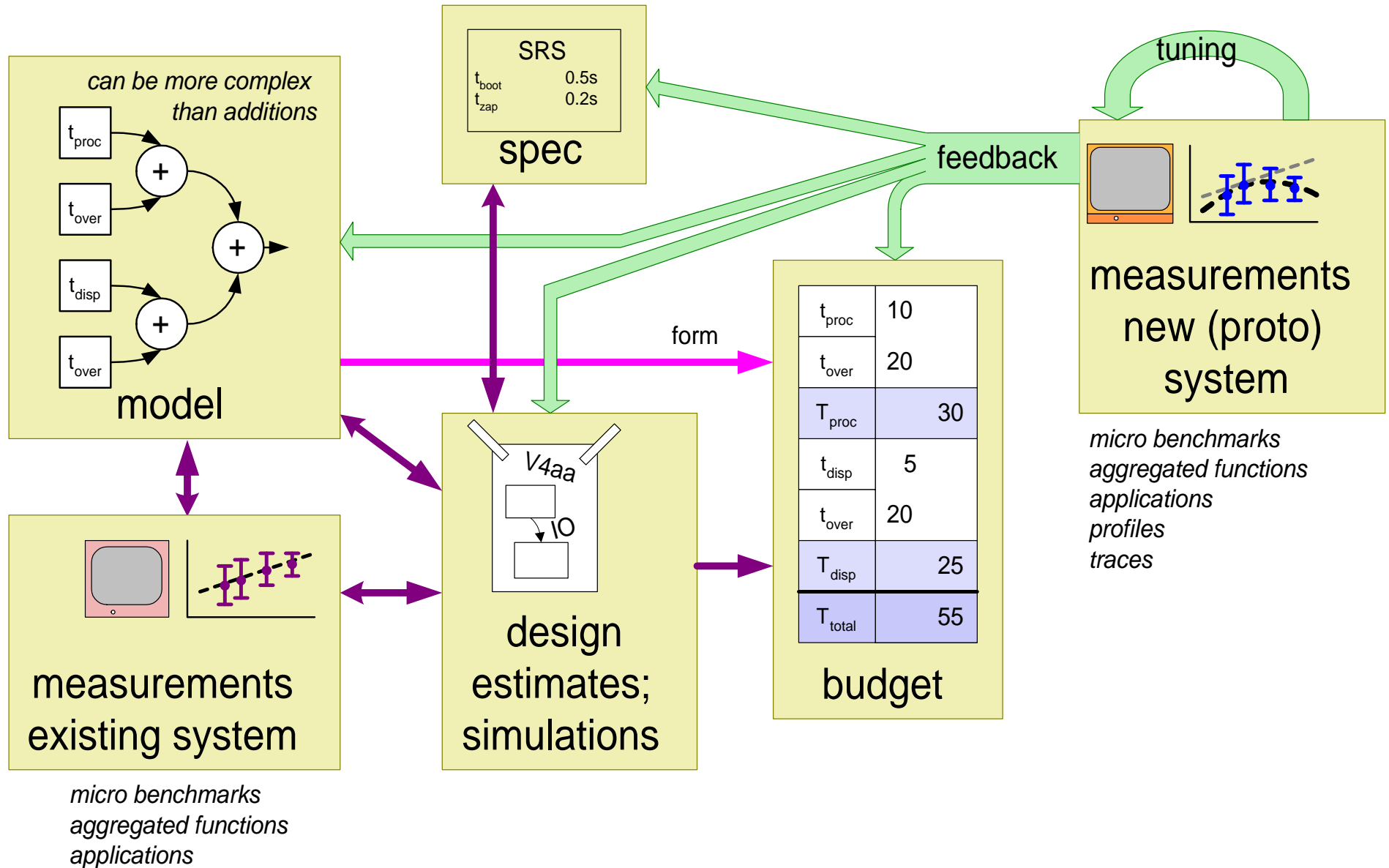
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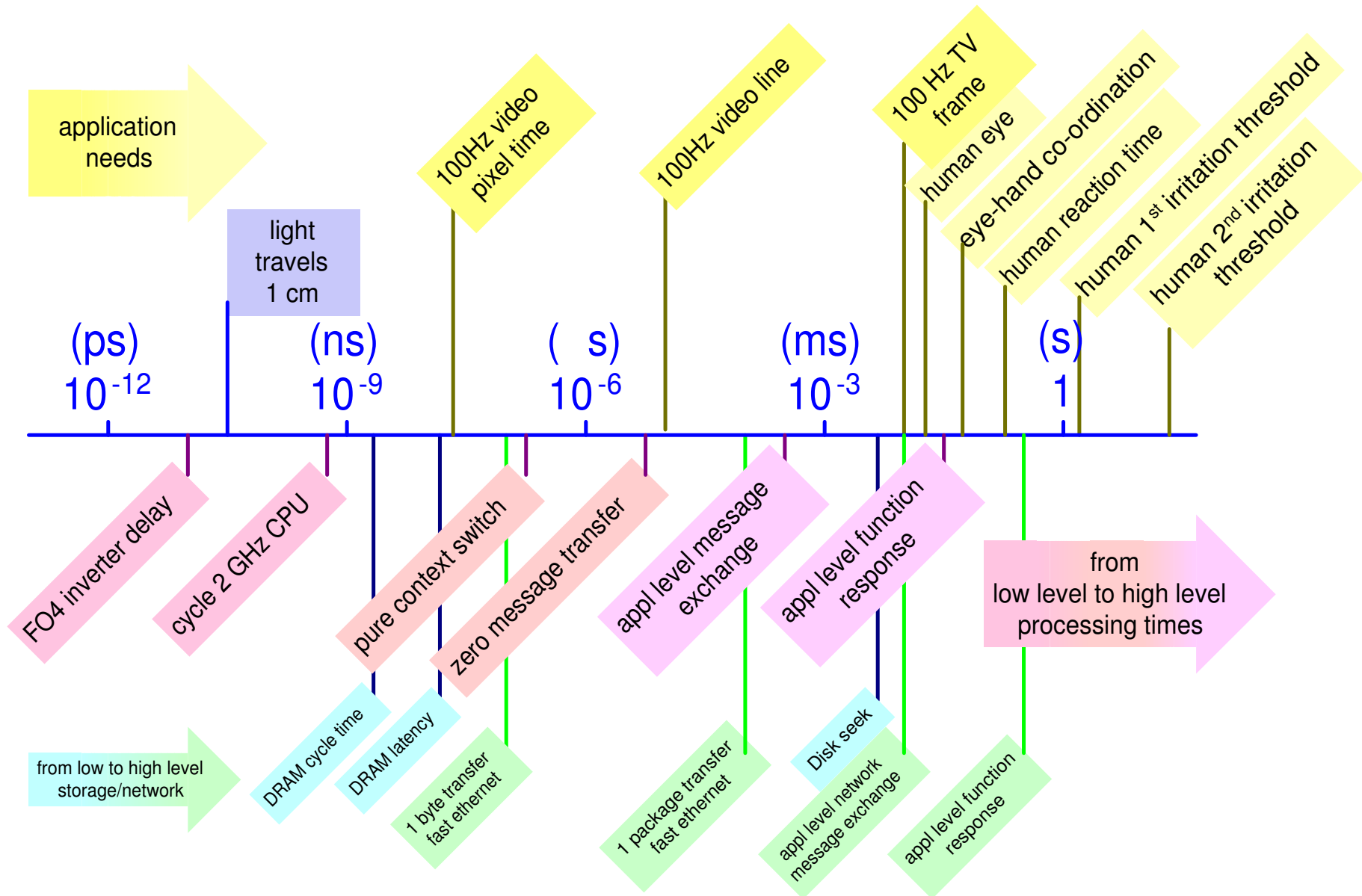
Budget based design flow



Example of a memory budget

<i>memory budget in Mbytes</i>	code	obj data	bulk data	total
shared code	11.0			11.0
User Interface process	0.3	3.0	12.0	15.3
database server	0.3	3.2	3.0	6.5
print server	0.3	1.2	9.0	10.5
optical storage server	0.3	2.0	1.0	3.3
communication server	0.3	2.0	4.0	6.3
UNIX commands	0.3	0.2	0	0.5
compute server	0.3	0.5	6.0	6.8
system monitor	0.3	0.5	0	0.8
application SW total	13.4	12.6	35.0	61.0
UNIX Solaris 2.x				10.0
file cache				3.0
total				74.0

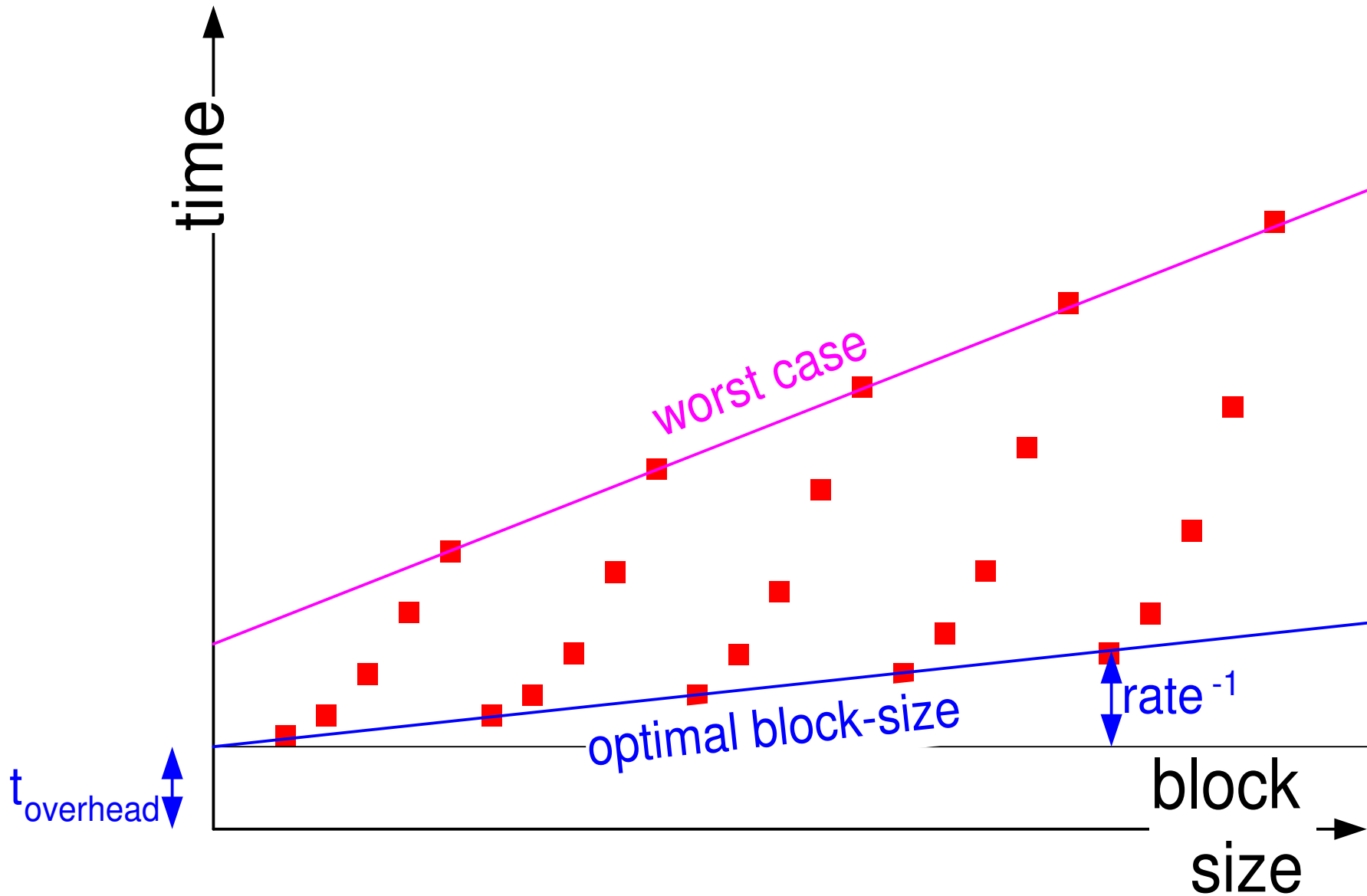
Actual timing on logarithmic scale



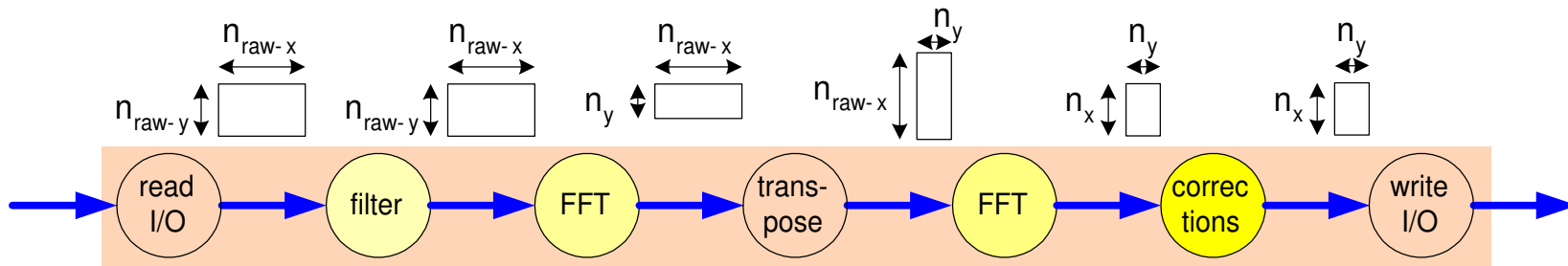
Typical micro benchmarks for timing aspects

	<i>infrequent operations, often time-intensive</i>	<i>often repeated operations</i>
<i>database</i>	start session finish session	perform transaction query
<i>network, I/O</i>	open connection close connection	transfer data
<i>high level construction</i>	component creation component destruction	method invocation same scope other context
<i>low level construction</i>	object creation object destruction	method invocation
<i>basic programming</i>	memory allocation memory free	function call loop overhead basic operations (add, mul, load, store)
<i>OS</i>	task, thread creation	task switch interrupt response
<i>HW</i>	power up, power down boot	cache flush low level data transfer

The transfer time as function of blocksize



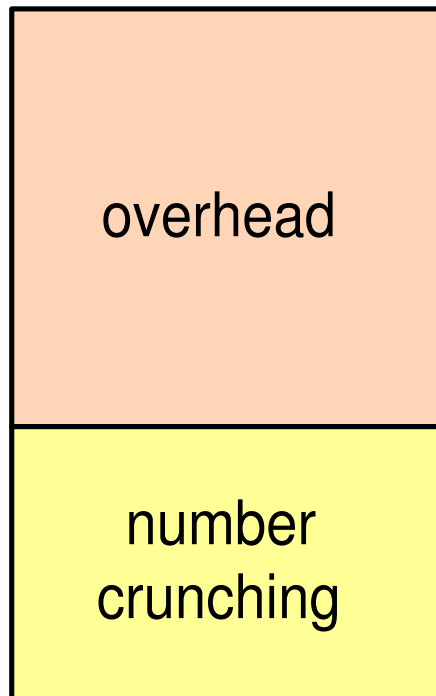
Performance evaluation



$$t_{\text{recon}} = t_{\text{filter}}(n_{\text{raw-x}}, n_{\text{raw-y}}) + n_{\text{raw-x}} * (t_{\text{fft}}(n_{\text{raw-y}}) + t_{\text{col-overhead}}) + n_y * (t_{\text{fft}}(n_{\text{raw-x}}) + t_{\text{row-overhead}}) + t_{\text{corrections}}(n_x, n_y) + t_{\text{read I/O}} + t_{\text{transpose}} + t_{\text{write I/O}} + t_{\text{control-overhead}}$$

$$t_{\text{fft}}(n) = c_{\text{fft}} * n * \log(n)$$

bookkeeping
transpose
malloc, free
write I/O
read I/O
overhead
correction computations
row overhead
FFT computations
column overhead
FFT computations
overhead
filter computations



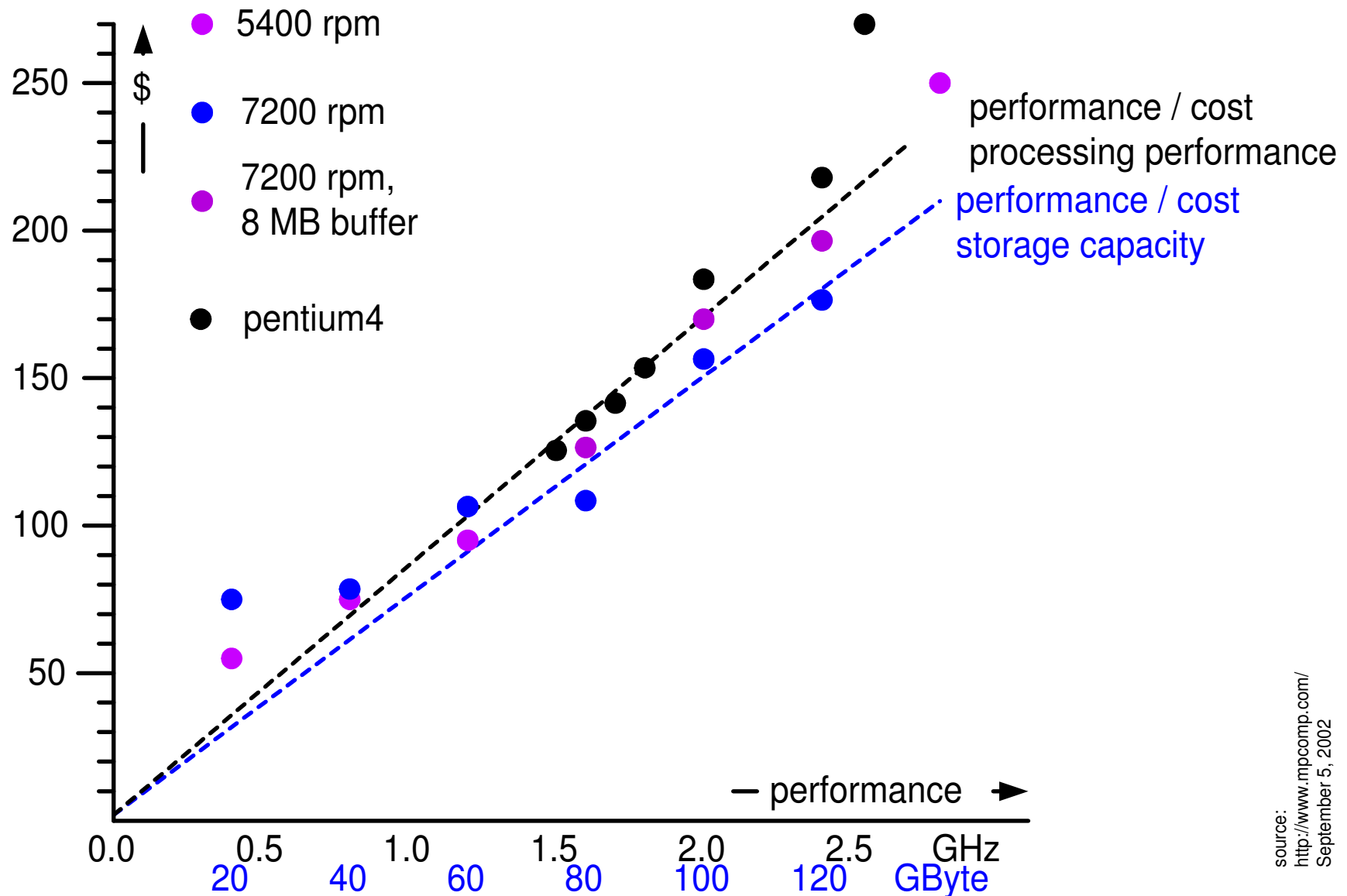
focus on overhead reduction

is more important

than faster algorithms

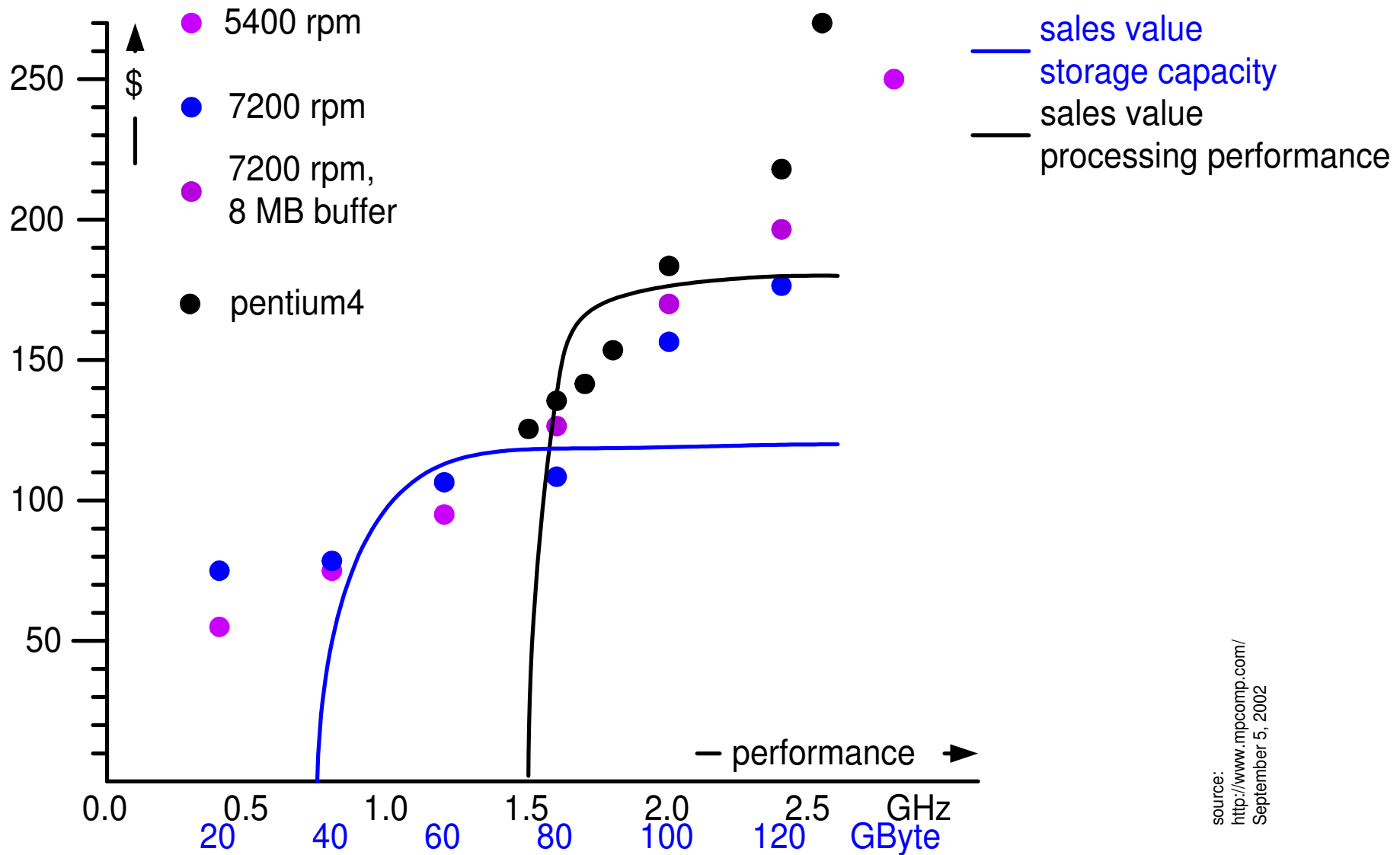
this is not an excuse for sloppy algorithms

Performance Cost, input data



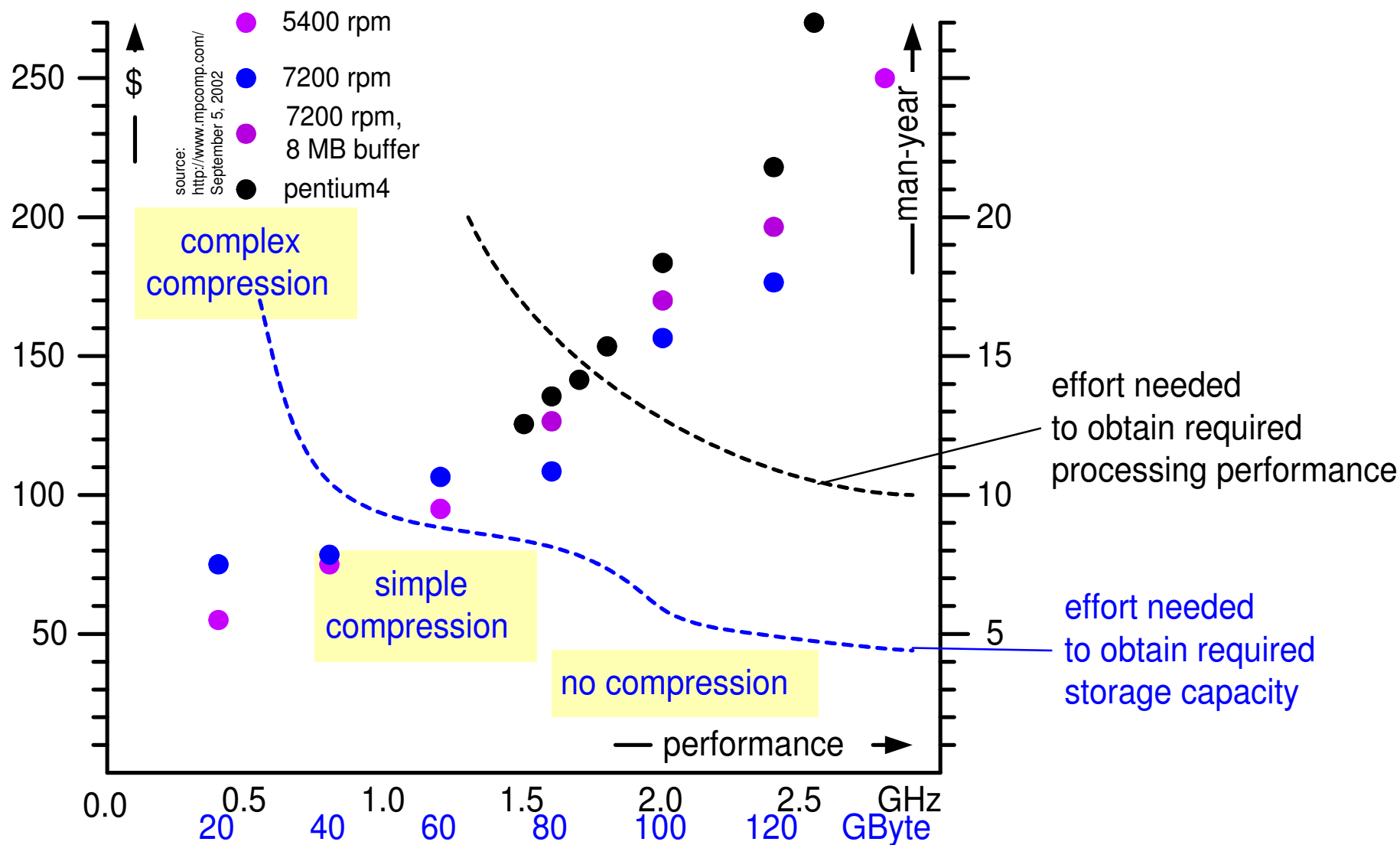
source: <http://www.mpcomp.com/>
September 5, 2002

Performance Cost, choice based on sales value

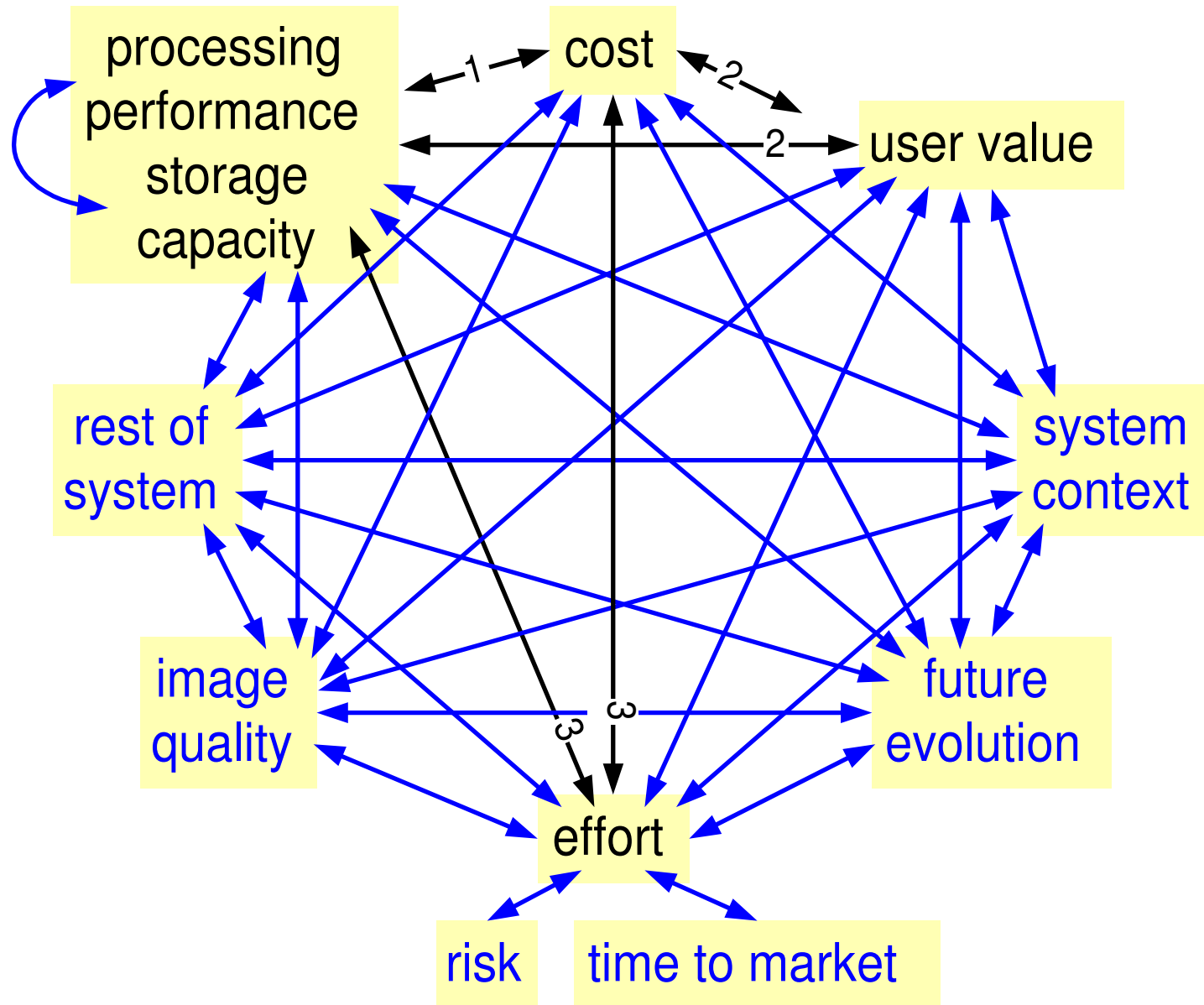


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<http://www.mpcomp.com/>
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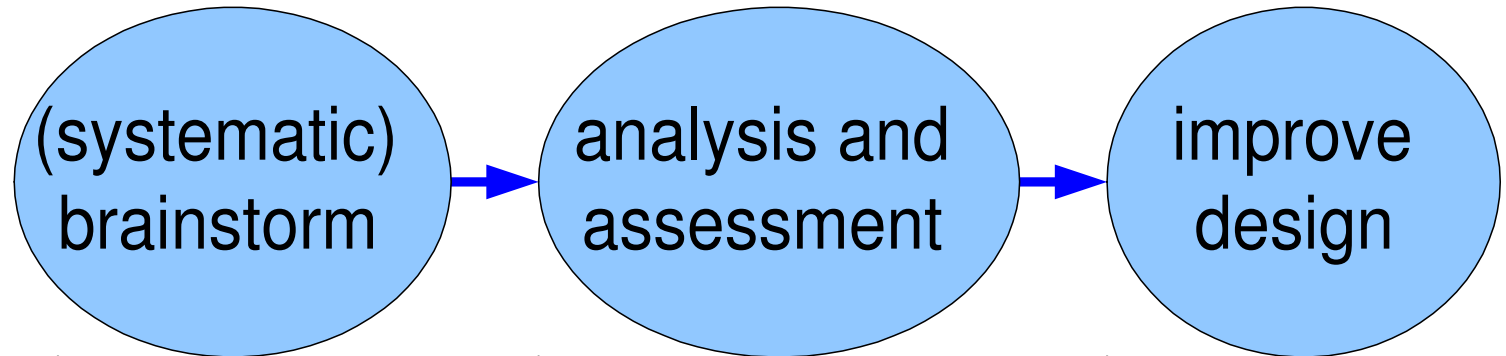
Performance Cost, effort consequences



But many many other considerations



Safety, Reliability and Security analysis methods



safety hazard analysis	potential hazards	probability severity	measures
reliability FMEA	failure modes	effects	measures
security	vulnerability risks	consequences	measures